Dynamic Adaptation of Replication Strategies for Reliable Agents

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Abstract

To make large-scale multi-agent systems reliable, we propose an adaptive application of replication strategies. Critical agents are replicated to avoid failures. As criticality of agents may evolve during the course of computation and problem solving, we need to dynamically and automatically adapt the number of replicas of agents, in order to maximize their reliability and availability based on available resources. We are studying an approach and mechanisms for evaluating the criticality of a given agent (based on application-level semantic information, e.g. messages intention, and also system-level statistical information, e.g., communication load) and for deciding what strategy to apply (e.g., active replication, passive) and how to parameterize it (e.g., number of replicas).
In this paper, we first present the replication mechanism and the framework named DarX that we developed to replicate agents. We then describe a new model to evaluate dynamically the criticality of agents. Then we describe the implementation of this model with the DarX fault-tolerant framework.

1 Introduction

A multi-agent system is a set of autonomous and interactive entities called agents (Avouris and Gasser, 1992). Recent real-life applications (e.g., intensive care monitoring, air traffic control and process control) are often distributed and must run continuously without any interruption. As a distributed system, however, multi-agent systems are exposed to possibility of failure of their hardware and/or software components. The failure of one component can often evolve into the failure of the whole system. To make these large-scale multi-agent systems reliable, an obvious solution is the introduction of redundancy: duplication (replication) of the critical components.

Replication mechanisms have been successfully applied for various distributed applications (Guerraoui and Schiper, 1997), e.g. data-bases. But in most cases, replication is decided by the programmer and applied statically, before the application starts. This works fine because the criticality of components (e.g., main servers) may be well identified and remains stable during the application run.

Opposite to that, in the case of dynamic and adaptive multi-agent applications, the criticality of agents may evolve dynamically during the course of computation. Moreover, the available resources are often limited. Thus, simultaneous replication of all the components of a large-scale system is not feasible. The idea is then to *automatically* and *dynamically* apply replication mechanisms *where* (to which agents) and *when* it is most needed. In this paper we will describe our approach to this objective and the realized tool to build easily reliable multi-agent systems.

This paper is organized as follows. Section 2 presents fault tolerance concepts and replication principles. Section 3 introduces a new approach of dynamic control of replication. Section 4 presents the DarX framework that we developed to replicate agents. This framework introduces novel features for dynamic control of replication. Section 5 describes our approach to compute agent criticality in order to guide replication. Section 6 describes the implementation of this solution and our preliminary experiments.

2 Fault-Tolerance

2.1 First and Simple Example

We consider the example of a distributed multi-agent system that helps at scheduling meetings. Each user has a personal assistant agent which manages his calendar. This agent interacts with:

- the user to receive his meeting requests and associated information (a title, a description, possible dates, participants, priority, etc.) ,
- the other agents of the system to schedule a meeting.

If the assistant agent of one important participant (initiator or prime participant) in a meeting fails (e.g., his ma-

chine crashes), this may disorganize the whole process. As the application is very dynamic - new meeting negotiations start and complete dynamically and simultaneously - decision for replication should be done automatically and dynamically.

2.2 Type of Faults

To achieve fault-tolerance, several distributed systems replicate their critical components. In this section, we define a classification of failures in multi-agent systems.

In computing systems, faults can for example occur in different sections of source code, and can result in a wide range of consequences. In distributed systems, another kind of faults can be considered: the communication failures.

A fault classification scheme is often used to categorize faults that have the same characteristics. The defined categories can then be used to collect statistics about faults and devise methods for fault prevention and detection. Thus, several classifications of failures have been proposed.

A. Fedoruk and R. Deters propose a classification of failures in multi-agent systems (Fedoruk and Deters, 2002). They group failures in five categories:

- Program bugs,
- Unforeseen states,
- Processor faults,
- Communication fault,
- Emerging unwanted behavior.

However, it is not easy to define the unwanted states of an agent because he is often adaptive. Moreover, a multi-agent system has no global control. So, the emerging unwanted behavior can be detected by an external observer, but it cannot be easily determined automatically by the system itself.

In our proposed solution, we consider, two categories of failures:

- Processor faults,
- Communication fault.

We think that our approach could also be applied to other categories of failures.

2.3 Techniques of Replication

This section, first, summarizes the principles of replication. It then points out the limits of current replication techniques and replication tools.

2.3.1 Principles of Replication

Replication of data and/or computation is an effective way to achieve fault tolerance in distributed systems. A replicated software component is defined as a software component that possesses a representation on two or more hosts (Guerraoui et al., 1989). There are two main types of replication protocols:

- active replication, in which all replicas process concurrently all input messages,
- passive replication, in which only one of the replicas processes all input messages and periodically transmits its current state to the other replicas in order to maintain consistency.

Active replication strategies provide fast recovery but lead to a high overhead. If the degree of replication is n, the n replicas are activated simultaneously to produce one result.

Passive replication minimizes processor utilization by activating redundant replicas only in case of failures. That is: if the active replica is found to be faulty, a new replica is elected among the set of passive ones and the execution is restarted from the last saved state. This technique requires less CPU resources than the active approach but it needs a checkpoint management which remains expensive in processing time and space.

The active replication provides a fast recovery delay. This kind of technique is dedicated to applications with real-time constraints which require short recovery delays. The passive replication scheme has a low overhead under failure free execution but does not provide short recovery delays. The choice of the most suitable strategy is directly dependent on the environment context, especially the failure rate, and the application requirements in terms of recovery delay and overhead. Active approaches should be chosen either when the failure rate becomes too high or the application design specifies hard time constraints. Otherwise, passive approaches are preferable.

2.3.2 Limits of Current Replication Techniques

Many toolkits (e.g., (Guerraoui et al., 1989) and (van Renesse et al., 1996)) include replication facilities to build reliable applications. However, most of them are not quite flexible enough to implement adaptive replication mechanism.

Therefore we designed a specific and novel framework for replication, named DarX (see details in section 4), which allows dynamic replication and dynamic adaptation of the replication policy (e.g., passive to active, changing the number of replicas).

3 Towards Dynamic Replication and Adaptive Control

Several solutions have been proposed to replicate distributed systems. These solutions are often used by the designer to replicate the system components before run time. The number of replicas and the replication strategy are explicitly and statically defined by the designer before run time. However, these solutions are not suitable to multi-agent systems. The solution we propose is mainly characterized by dynamic replication and adaptive control.

3.1 Dynamic Replication

Several replication strategies (mainly, active and passive) can be used to replicate Agents. As explained in Section 2.3.1, each strategy has its pros and cons, the tradeoff being recovery speed versus overhead. Thus, the choice of the most suitable strategy relies on the environment context.

In most multi-agent applications, the environment context is very dynamic. So, the choice of the replication strategy of each component, which relies on a part of this environment, must be determined dynamically and adapted to the environment changes.

Moreover, a multi-agent system component which can be very critical at a moment can loose its critically later. If we consider the replication cost which is very high, the number of replicas of these components must be therefore dynamically updated.

Thus, the solution we propose allows to dynamically adapt the number of replicas and the replication strategy. This solution is provided by the framework DARX (see section 4).

3.2 Adaptive Control

DarX provides the needed adaptive mechanisms to replicate agents and to modify the replication strategy. Meanwhile, we cannot always replicate all the agents of the system because the available resources are usually limited. In the given example (section 2.1), we can consider more than 100 assistant agents and resources that do not allow to duplicate more than 60 agents. The problem therefore is to determine the most critical agents and then the needed number of replicas of these agents.

We distinguish two cases:

- the agent's criticality is static,
- the agent's criticality is dynamic.

In the first case, multi-agent systems have static organization structures, static behaviors of agents, and a small number of agents. Critical agents can be therefore identified by the designer and can be replicated by the programmer before run time.

In the second case, multi-agent systems may have dynamic organization structures, dynamic behaviors of agents, and a large number of agents. So, the agents criticality cannot be determined before run time. The agent criticality can be therefore based on these dynamic organizational structures. The problem is how to determine dynamically these structures to evaluate the agent criticality? Thus, we propose a new approach for observing the domain agents and evaluating dynamically their criticality. This approach is based on two kinds of information: semantic-level information and system-level information.

4 Darx

DarX is a framework to design reliable distributed applications which include a set of distributed communicating entities (agents). Each agent can be dynamically replicated an unlimited number of times and with different replication strategies (passive and active).

4.1 Darx Architecture

DarX includes group membership management to dynamically add or remove replicas. It also provides atomic and ordered multi-cast for the replication groups' internal communication. For portability and compatibility issues, DarX is implemented in Java.

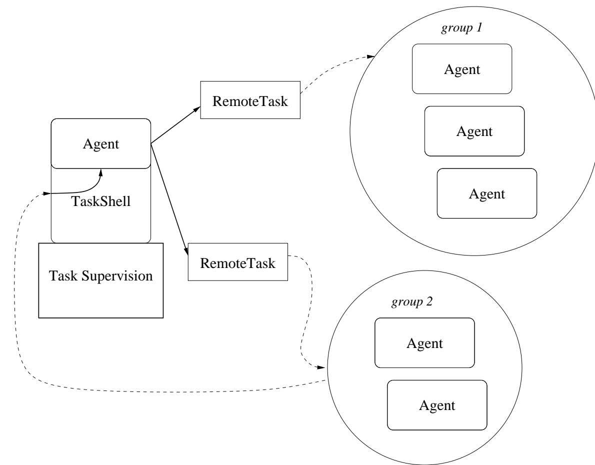

Figure 1: DarX application architecture

A replication group is an opaque entity underlying every application agent. The number of replicas and the internal strategy of a specific agent are totally hidden to the other application agents. Each replication group has exactly one leader which communicates with the other agents. The leader also checks the liveness of each replica and is responsible for reliable broadcasting. In case of failure of a leader, a new one is automatically elected among the set of remaining replicas.

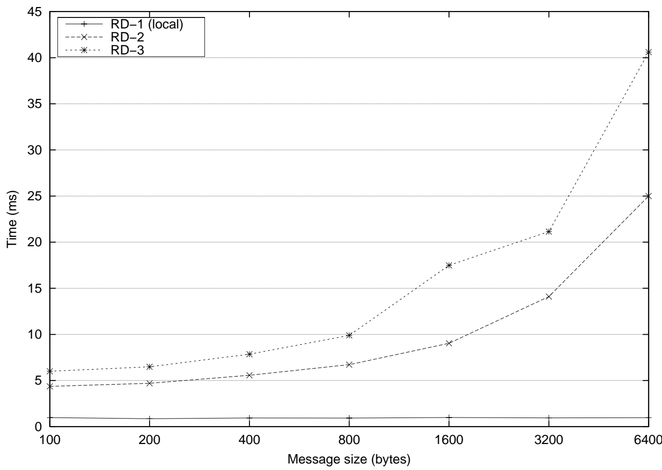

Figure 2: Communication cost as a function of the replication degree

DarX provides global naming. Each replicated agent has a global name which is independent of the current location of its replicas. The underlying system allows to handle the agent's execution and communication. Each agent is itself wrapped into a TaskShell (Figure 1), which acts as a replication group manager and is responsible for delivering received messages to all the members of the replication group, thus preserving the transparency for the supported application. Input messages are intercepted by the TaskShell, enabling message caching. Hence all messages get to be processed in the same order within a replication group.

An agent can communicate with a remote agent, unregarding whether it is a single agent or a replication group, by using a local proxy implemented by the RemoteTask interface. Each RemoteTask references a distinct remote entity considered as the leader of its replication group. The reliability features are thus brought to agents by an instance of a DarX server (DarxServer) running on every location. Each DarxServer implements the required replication services, backed up by a common global naming/location service.

4.2 Measurements

Our first experiments and measurements of DarX are very promising. We evaluated several costs and made comparisons with other systems (see (Marin et al., 2001)).

In this paper, we just show the cost of sending a message to a replication group using the active replication strategy. Figure 2 presents three configurations with different replication degrees. In the RD-1 configuration, the task is local and not replicated. In the RD-2 (resp. RD-3) configuration, there is one (resp. two) replica(s); the leader being on the sending host and the other replica(s) residing on one (or two) distinct remote host(s).

5 Adaptive Control of Replication

We will now detail our approach for dynamically evaluating criticality of each agent in order to perform dynamic replication where and when best needed.

5.1 Hypothesis and principles

We want some automatic mechanism for generality reasons. But in order to be efficient, we also need some prior input from the designer of the application. This designer can choose among several approaches of replication: static and dynamic.

In the proposed dynamic approach, the agent criticality relies on two kinds of information:

- System-level information. It will be based on standard measurements (communication load, processing time...). We are currently evaluating their significance to measure the activity of an agent.
- Semantic-level information.

Several aspects may be considered (importance of agents, independence of agents, importance of messages...). We decided to use the concept of role, because it captures the importance of an agent in an organization, and its dependencies to other agents.

Note that our approach is generic and that it is not related to a specific interaction language or application domain. Also agents can be either reactive or cognitive. We just suppose that they communicate with some agent communication language such as ACL (FIPA, 1997) and KQML (Finin et al., 1994).

5.2 Example

The application designer will manually evaluate criticality of the roles, corresponding to their ”importance” in the organization and in the computation.

In the example introduced in section 2.1, we are considering two roles: Initiator and Participant (Finin et al., 1994). Their respective weights will be set by the application designer to respectively 0.7 and 0.3 (see 1).

Table 1: Examples of roles and their weights

Roles	Weights
Initiator	0.7
Participant	0.3

5.3 Architecture

In order to track the dynamical adoption of roles by agents, we propose a role recognition method. Our approach is based on the observation of the agent execution and their

interactions to recognize the roles of each agent and to evaluate his processing activity. This is used to dynamically compute the criticality of an agent.

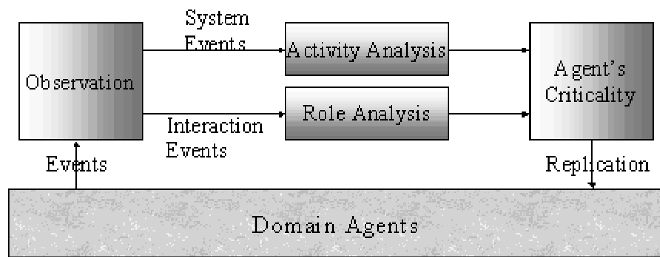

Figure 3: General architecture for replication control

In order to collect the data, we associate an observation module to each DarxServer on each machine (see section 4). This module will collect events (provided by DarxServer). A role analysis is then associated to each agent (leader of the group replica) of this machine, by considering his sent and received messages.

The basic architecture controlling the replication of agents is shown in Figure 3.

The next sections describe the role analysis and activity analysis methods that we propose.

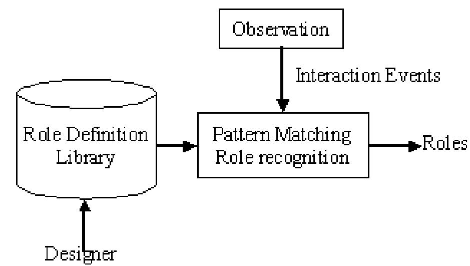

Figure 4: Roles recognition

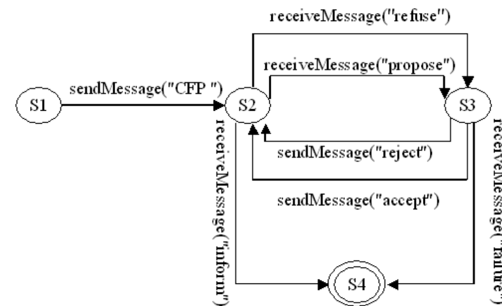

Figure 5: Example of ATN

5.4 Role Analysis

We consider two cases. In the first case, each agent displays explicitly his roles or interaction protocols. The roles of each agent are thus easily deduced from interaction events. In the second case, agents do not display their roles nor their interaction protocols. The agent roles are deduced from the interaction events by the role analysis module.

In this analysis, attention is focused on the precise ordering of interaction events (exchanged messages). The role module captures and represents the set of interaction events resulting from the domain agent interactions (sent and received messages).

We associate to each agent an entity that analyses the associated interaction events. This analysis determines the roles of the agent. Figure 4 illustrates the various steps of this analysis.

To represent the agent interactions, several methods have been proposed such as state machines and Petri nets (Fallah-Seghrouchni et al., 1999). For our application, state machines provide a well suitable representation. Each role interaction model is represented by an augmented transition network (ATN) (Woods, 1970). A transition represents an interaction event (sending or receiving a message). Figure 5 shows an example of ATN that represents the interaction model of the role Initiator described below.

A library of roles definition is used to recognize the active roles. To facilitate the initialization of this library, we have introduced a role description language. Each role is represented by a set of interaction events. This language is based on a set of operators (similar to those proposed in (M. Wooldridge and Kinny, 1999), see Table 2), interaction events and variables.

Interaction events represent the exchanged messages. We distinguish two kinds of interaction events: ReceiveMessage and SendMessage. The attributes of the SendMessage and ReceiveMessage interaction events are similar to the attributes of ACL messages:

- SendMessage(Communicative act, sender, receiver, content, reply-with, ...).
- ReceiveMessage(Communicative act, sender, receiver, content, reply-with, ...).

Table 2: Operators

Operators	Interpretation
A.B	Separate two consecutive events
A|B	Or
A||B	Parallel events
(A)*	O time or more
(A)+	1 time or more
(A)n	n time or more
[A]	Facultative

In order to be able to filter various messages, we introduce the "wild card" character ?. For example, in the interaction event ReceiveMessage ("CFP", "X", "Y", ?), the content is unconstrained. So, this interaction event can match any other interaction event with the communication act CFP, the sender "X", the receiver "Y" and any contents.

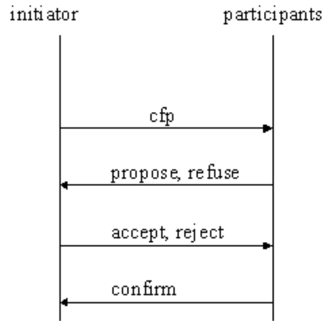

Figure 6: Contract net protocol

In the example of scheduling meetings, the assistant agents use the contract net protocol (FIPA, 1997) (see 6) to schedule a meeting. The interaction model of the initiator role is deduced from the contract net protocol. It is described in Table 3.

This description represents the different steps (sent and received messages) of the Initiator. It can be interpreted as follows (FIPA, 1997).

- A call for proposals message is sent to the participants from the initiator following the FIPA Contract Net protocol.
- The participants reply to the initiator with the proposed meeting times. The form of this message is either a proposal or a refusal.
- The initiator sends accept or reject messages to participants.
- The participants which agree to the proposed meeting inform the initiator that they have completed the request to schedule a meeting (confirm).

Table 3: Description of the role *Initiator*

(SendMessage("CFP", Agent,?,?,M1))+ . ((ReceiveMessageEvent("propose", ?, Agent,?,M2, M1)) | (ReceiveMessageEvent("refuse", ?, Agent, ?,M2, M1)))+ . ((SendMessage("accept", Agent,?,?, M2, M1)) | (SendMessage("reject", Agent,?,?, M2, M1)))+ . (ReceiveMessageEvent("confirm", ?, Agent,?,M2))+ .

Note that in many cases, roles can be deduced before the end of the associated sequence of interaction events (final state of the associated ATN). In the scheduling meetings example, the role Initiator may be recognized as soon as the "CFP" message is received, as it is unique to this role.

5.5 Activity Analysis

In multi-agent systems, the internal activity of agents cannot be observed, because it is private. The observation is restricted to events. To evaluate the degree of the agent activity, we use system events that are collected at the system level. We are considering two kinds of events: CPU time and communication load. We are currently evaluating the significance of these measures as indicators of agent activity, to be useful to calculate agent criticality.

For an agent $Agent_i$ and a given time interval Δt, these events provide:

- The used time of CPU (cp_i),
- The communication load (cl_i).

cp_i and cl_i may be then used to measure the agent degree of activity aw_i as follows:

$$aw_i = (d_1 * cp_i/\Delta t + d_2 * cl_i/CL)/(d_1 + d_2) \quad (1)$$

where:

- CL is the global communication load,
- d_1 and d_2 are weights introduced by the user.

5.6 Agent Criticality

The analysis of events (system events and interaction events) provides two kinds of information: the roles and the degree of activity of each agent. This information is then processed by the agent's criticality module. The latter relies on a table T (an example is given in Table 1) that defines the weights of roles. This table is initialized by the application designer. Table 3 gives examples of roles and their weights.

The criticality of the agent $Agent_i$ which fulfills the roles r_{i1} to r_{im} is computed as follows:

$$w_i = (a_1 * \sum_{j=1,m} T[r_{ij]}] + a_2 * aw_i)/(a1 + a2) \quad (2)$$

Where: a_1 and a_2 are the weights given to the two kinds of parameters (roles and degree of activity). They are introduced by the designer.

For each Agent A_i, its criticality w_i is used to compute the number of his replicas.

5.7 Replication

An agent is replicated according to:

- w_i: his criticality,
- W: the sum of the domain agents' criticality,
- rm: the minimum number of replicas which is introduced by the designer,
- Rm: the available resources which define the maximum number of possible simultaneous replicas.

The number of replicas nb_i of $Agent_i$ can be determined as follows:

$$nb_i = rounded(rm + w_i * Rm/W) \quad (3)$$

Table 4: Examples of agents, their weights and the associated number of replicas

Agents	Criticality per agent	Number of replicas per agent
Agent1, Agent2, Agent3, Agent 4	0,9	2
Agent5, Agent6, Agent7, Agent8, Agent9, Agent10, Agent11, Agent12, Agent13, Agent14	0.5	1
Agent15, Agent16, Agent17, Agent18, Agent19, Agent20, Agent21, Agent22, Agent23, Agent24	0.2	0

Table 4 gives an example of agents, their criticality and the associated replicas when Rm = 20 and rm= 0. Note that (rm=0) means that the agent is not replicated.

The numbers of replicas are then used by DarX to update the number of replicas of each agent.

6 Experiments

We made some preliminary experiments using the scenario of agents scheduling their meetings, as introduced in section 2.1.

Agents take randomly roles of Initiator, choose Participants for scheduling meetings or remain inactive (without any role). Several meetings are scheduled simultaneously. The number of critical agents (which can be either Initiator or Participant) is 60% of the number of agents.

As the distributed observation module implementation has not been completed yet, we have run these preliminary experiments on a single machine. In order to simulate the presence of faults, we implemented a failure simulator randomly stopping the thread of an agent (chosen randomly). The number of the introduced faults was set equal to the number of agents. We repeated several times the experiments with a variable number of resources (number of replicas that can be used).

From these first experiments, we found that the number of resources should be at least equal to the number of critical agents.

We are currently working on more experiments and measurements in order to better evaluate our adaptive control architecture and to compare it to other control methods (including random replication).

7 Related Work

Several approaches address the multi-faced problem of fault tolerance in multi-agent systems. These works can be classified in two main categories. A first approach focuses especially on the reliability of an agent within a multi-agent system. This approach handles the serious problems of communication, interaction and coordination of agents (and their replicas) with the other agents of the system. The second approach addresses the difficulties of making reliable an agent, particularly a mobile agent, which is more exposed to security problems (Pleisch and Schiper, 2001) (Silva and Popescu-Zeletin, 1998) (Johansen et al., 199) (Strasser et al., 1998). This second approach is beyond the scope of this paper.

Within the family of reactive multi-agent systems, some systems offer high redundancy. A good example is a system based on the metaphor of ant nests. Unfortunately:

- we cannot design any application in term of such reactive multi-agent systems. Basically we do not have yet a good methodology.
- we cannot apply such simple redundancy scheme onto more cognitive multi-agent systems as this would cause inconsistencies between copies of a single agent.

Some work (Decker et al., 1997) offers dynamic cloning of specific agents in multi-agent systems. But their motivation is different, the objective is to improve the availability of an agent if it is too congested. The agents considered seem to have only functional tasks (with no changing state) and fault-tolerance aspects are not considered.

S. Hagg introduces sentinels to protect the agents from some undesirable states (Hagg, 1997). Sentinels represent the control structure of their multi-agent system. They need to build models of each agent to perform functionalities and monitor communications in order to react to faults. Each sentinel is associated by the designer to one functionality of the multi-agent system. This sentinel handles the different agents which interact to achieve the functionality. The analysis of his believes on these agents enables the sentinel to detect a fault when it occurs. Adding

sentinels to multi-agent systems seems to be a good approach, however the sentinels themselves represent failure points for the multi-agent system.

(Kumar et al., 2000) present a fault tolerant multi-agent architecture that regroups agents and brokers. They address the problem of recovering the multi-agent system from only its broker failures.

(Fedoruk and Deters, 2002) propose to use proxies. This approach tries to make transparent the use of agent replication, i.e. enabling the replicas of an agent to act as a same entity regarding the other agents of the system which will not know that they are interacting with a group of replicas. The proxy manages the state of the replicas. To do so, all the external and internal communications of the group are redirected to the proxy. However this increases the workload of the proxy which is a quasi central entity. To make it reliable, they propose to build, for instance, a hierarchy of proxies for each group of replicas. They point out the specific problems of read/write consistency, resource locking also discussed in (Silva et al., 2000).

In distributed computing, many toolkits include replication facilities to build reliable application. However, many of products are not enough flexible to implement an adapted replication. MetaXa (M.Golm, 1998) implements in Java active and passive replication in a flexible way. Authors extended Java with a reactive metalevel architecture. Like in DarX, the replication is transparent. However, MetaXa relies on a modified Java interpreter. GARF (Guerraoui et al., 1989) realizes fault-tolerant Smalltalk machines using active replication. Similar to MetaXa, GARX uses a metalevel and provides different replication strategies. But, it does not provide adaptive mechanism to apply these strategies.

8 Conclusion

Large-scale multi-agent systems are often distributed and must run without any interruption. To make these systems reliable, we proposed a new approach to evaluate dynamically the criticality of agents. This approach is based on the concepts of roles and degree of activity. The agent criticality is then used to replicate agents in order to maximize their reliability and availability based on available resources.

To validate the proposed approach, we realized a fault-tolerant framework (Darx). The integration of DARX with a multi-agent platform, such as DIMA ieee99, provides a generic fault-tolerant multi-agent platform. In order to validate this fault-tolerant multi-agent platform, two small applications have been developed (meetings scheduling and crisis management system). They are intended at evaluating our model and architecture viability. They aim also at completing the model and adjusting the parameters. The obtained results are interesting and promising. However, more experiments with real-life applications are needed to validate the proposed approach.

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